

own copy of the document to note which of the collaborators have acknowledged receipt of the document.

[0190] Arrows 1745 and 1750 illustrate a logout scenario. Before broadcasting a LOGOUT message, the computer 252 constructs the LOGOUT message from the predetermined LOGOUT opcode, the username of the collaborator logging out (i.e., Z) of a current editing session, and the UUID of the document. Then the computer 252 being operated by Z transmits the LOGOUT message to the computer 200, as represented by the arrow 1745. The computer 252 also transmits the LOGOUT message to the computer 250 being operated by Y, as reference by the arrow 1745.

[0191] FIG. 18 is a sequence diagram showing the process of performing an online merge of two versions of a document with both simultaneous and sequential modifications being made to the document. Again, the collaborator X is operating the computer 200 and the collaborator Y is operating the computer 250. The arrows referenced as 1810, 1815, 1820, and 1825 show a scenario where the collaborators X and Y make simultaneous modifications while online. The computer 200 constructs a MODIFY message from the predetermined MODIFY opcode, the username (X) of the modifier, the UUID of the document, the old version string (0), the new version string (1), and the difference between the two versions of the document. The computer 250 also constructs a MODIFY message from the predetermined MODIFY opcode, from the predetermined MODIFY opcode, the username (Y) of the modifier, the UUID of the document, the old version string (0), the new version string (1), and the patch.

[0192] Each copy of the document preferably contains a history of modifications to the document and a record of who made the modifications and, insofar as the document is aware, which collaborators have been made aware of each of the modifications. Accordingly, as described above, associated with each document is version information comprising a version identifier or version string and a list of "known acknowledgers", one of whom is the author. This is represented in FIGS. 17, 18, and 19 by the notation doc(version identifier, known acknowledgers), e.g. doc(0, {X, Y}) to indicate that as far as a particular collaborator is aware, version 0 has been acknowledged by X and Y, but not by Z.

[0193] The computer 200 transmits the constructed MODIFY message to the computer 250 as represented by the arrow 1810. The MODIFY message includes the difference between the two versions of the document. Then the computer 250 being operated by Y transmits the MODIFY message which has been constructed by the computer 250 to the computer 200. This MODIFY message includes the patch representing the changes made to the document by Y during the editing session.

[0194] The computers 200 and 250 being operated by X and Y must then construct acknowledgment messages for the modifications that they have sent each other. The computer 200 constructs a modification acknowledgment message, ACK_MOD, from the predetermined ACK_MOD opcode. The ACK_MOD message includes the modifier username (i.e., Y), the username of the acknowledger (i.e., X), the UUID of the document, the version string of the version to which the patch applied, and, optionally, a hash of the patch included with the MODIFY message sent by Y (i.e., for verification purposes). The computer 250 constructs

a modification acknowledgment message, ACK_MOD, from the predetermined ACK_MOD opcode. Again, the ACK_MOD message includes the username of the modifier (i.e., X), the username of the acknowledger (i.e., Y), the document UUID, the version string of the version to which the patch applied, and, optionally, a hash of the patch representing the modifications to the document made by X (i.e., again for verification purposes). As represented by the arrow 1820, the computer 250 being operated by Y transmits the ACK_MOD message to computer 200. Then, the computer 200 being operated by X transmits the ACK_MOD message to computer 250, as represented by the arrow 1825.

[0195] The arrows 1830 and 1835 show the sequence where user X modifies the document, the computer 200 constructs and transmits a MODIFY message, as described above, and the computer 250 being operated by the user Y acknowledges the modification by constructing and transmitting an ACK_MOD message to computer 200. Similarly, the arrows 1840 and 1845 show the sequence where user Y modifies the document, the computer 250 constructs and transmits a MODIFY message, as described above, and the computer 200 being operated by the user X acknowledges the modification by constructing and transmitting an ACK_MOD message to computer 250 which will update the copy of the document to record X's acknowledgement.

[0196] Note that throughout FIG. 18 no messages are transmitted to collaborator Z. This is because collaborators X and Y removed collaborator Z from their respective lists of active collaborators in response to the LOGOUT message broadcast by collaborator Z at steps 1745 and 1750.

[0197] FIG. 19 is a sequence diagram showing the sequence of events for an online merge of two versions of a document with both simultaneous and sequential modifications. Again, the collaborator X is operating the computer 200 and the collaborator Y is operating the computer 250. Then, the collaborator Z logs into an editing session using the computer 252. The arrows 1915, 1920, 1925, 1930, 1935, 1940, 1945, 1950, 1955, 1960, 1965, 1970, 1975, and 1980 show sequence of event when user Z logs back into the editing session.

[0198] Before attempting to log in, the computer 252 must construct a login message, LOGIN, from the predetermined LOGIN opcode, the username (i.e., Z), and the document UUID. The computer 252 then transmits the LOGIN message to the computer 200, as represented by the arrow 1915. The computer 252 also transmits the LOGIN message to the computer 250, as represented by the arrow 1920.

[0199] The computers 200 and 250 then acknowledge the LOGIN message transmitted by the computer 252 so that Z can discover which collaborators are online and can provide an update. As illustrated in FIG. 6, the computer 200 constructs an ACK_LOGIN message from the predetermined ACK_LOGIN opcode, the username (X), and the document's UUID, and then transmits the message to the computer 252, as represented by the arrow 1925. Similarly, the computer 250 constructs an ACK_LOGIN message from the predetermined ACK_LOGIN opcode, the username (Y), and the document's UUID, and then transmits the message to the computer 252, as represented by the arrow 1930.

[0200] The computer 252 being operated by Z must now update the version of the document stored locally on the